

Axiomatic Language

A Short Presentation

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Intl. Conf. On Logic Programming
Las Cruces, New Mexico
September, 2019

<http://axiomaticlanguage.org/>

Language Goals

1. pure specification – what, not how
2. minimal, but extensible
3. metalanguage – can subsume other languages

Specification by Enumeration

Idea: Program external behavior defined by infinite set of symbolic expressions that enumerate inputs and corresponding outputs.

Recipe

- pure, definite Prolog with Lisp syntax
- higher-order generalization [HiLog]
- string variables

The Core Language (1)

Axioms generate valid expressions.

expression:

atom, – `abc, `+

expression variable, – %x, %1

sequence of ≥ 0 expressions and string variables

– (`M () \$1 %)

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axiom: conclusion expr and ≥ 0 **condition** exprs

<conclu> < *<cond1>*, ..., *<condn>*.

<conclu>. ! unconditional axiom

The Core Language (2)

axiom instance: substitute values for variables in axiom

$(\text{`A } \%x \$1) < (\text{`B } \%x), (\text{`C } \$1).$

$\rightarrow (\text{`A } \text{`x } \text{`u } ()) < (\text{`B } \text{`x}), (\text{`C } \text{`u } ()).$

The Core Language (2)

axiom instance: substitute values for variables in axiom

$(\text{`A } \%x \$1) < (\text{`B } \%x), (\text{`C } \$1).$

→ $(\text{`A } \text{`x } \text{`u } ()) < (\text{`B } \text{`x}), (\text{`C } \text{`u } ()).$

valid expressions: If all conditions of an axiom instance are valid expressions, the conclusion is valid.

$(\text{`a } \text{`b}).$

$((\%) \$ \$) < (\% \$).$

→ $(\text{`a } \text{`b}),$

$((\text{`a}) \text{`b } \text{`b}),$

$(((\text{`a})) \text{`b } \text{`b } \text{`b } \text{`b}),$

...

Example – Natural Numbers

Set of natural numbers:

$(\text{`num } (\text{`z})) .$

$(\text{`num } (\text{`s } \$)) < (\text{`num } (\$)) .$

$\rightarrow (\text{`num } (\text{`s } \text{`s } \text{`z}))$

Addition of natural numbers:

$(\text{`plus } \%n (\text{`z}) \%n) < (\text{`num } \%n) .$

$(\text{`plus } \%1 (\text{`s } \$2) (\text{`s } \$3)) <$

$(\text{`plus } \%1 (\$2) (\$3)) .$

$\rightarrow (\text{`plus } (\text{`s } \text{`z}) (\text{`s } \text{`z}) (\text{`s } \text{`s } \text{`z}))$

Syntax Extensions

single char in single quotes:

```
'A' = ( `char ( `0 `1 `0 `0 `0 `0 `0 `1 ) )
```

char string in single quotes within sequence:

```
( ... 'abc' ... ) = ( ... 'a' 'b' 'c' ... )
```

char string in double quotes:

```
"abc" = ( 'abc' ) = ( 'a' 'b' 'c' )
```

symbol not starting with special char:

```
abc = ( ` "abc" )
```

Example – List Predicates

Concatenation of sequences:

```
(concat ($1) ($2) ($1 $2)).
```

```
→ (concat (a b) (c d e) (a b c d e))
```

Member of a sequence:

```
(member % ($1 % $2)).
```

```
→ (member b (a b c))
```

Reverse of a sequence:

```
(reverse () ()).
```

```
(reverse (% $seq) ($rev %))<
```

```
(reverse ($seq) ($rev)).
```

```
→ (reverse (u v) (v u))
```

Conclusion

- Specifications – software engineering benefit
 - Smaller, more readable, more reusable, more correct
- Minimal & pure – well-suited to proof
 - Equivalence of specification and program
 - Prove assertions to validate specification
- Billion-dollar application!
 - http://axiomaticlanguage.org/A_Vision_for_CAD_released.pdf
- Implementation grand challenge
 - Transformation of specifications to programs
 - http://axiomaticlanguage.org/LOPSTR18_LM_released.pdf